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Ian June L. Garces
Ateneo de Manila University, ijlgarces@ateneo.edu

Jose B. Rosario

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Computing the Metric Dimension of Truncated Wheels

Ian June L. Garces

Department of Mathematics
Ateneo de Manila University
1108 Loyola Heights, Quezon City
The Philippines

Jose B. Rosario

Department of Mathematics
Ateneo de Manila University
1108 Loyola Heights, Quezon City
The Philippines

Abstract

For an ordered subset $W = \{w_1, w_2, w_3, \ldots, w_k\}$ of vertices in a connected graph G and a vertex v of G, the metric representation of v with respect to W is the k-vector $r(v|W) = (d(v, w_1), d(v, w_2), d(v, w_3), \ldots, d(v, w_k))$. The set W is called a resolving set of G if $r(u|W) = r(v|W)$ implies $u = v$. The metric dimension of G, denoted by $\beta(G)$, is the minimum cardinality of a resolving set of G.

Let $n \geq 3$ be an integer. A truncated wheel, denoted by TW_n, is the graph with vertex set $V(TW_n) = \{a\} \cup B \cup C$, where $B = \{b_i : 1 \leq i \leq n\}$ and $C = \{c_{j,k} : 1 \leq j \leq n, 1 \leq k \leq 2\}$, and edge set $E(TW_n) = \{ab_i : 1 \leq i \leq n\} \cup \{b_i c_{i,k} : 1 \leq i \leq n, 1 \leq k \leq 2\} \cup \{c_{j,1}c_{j,2} : 1 \leq j \leq n\} \cup \{c_{j,2}c_{j+1,1} : 1 \leq j \leq n\}$, where $c_{n+1,1} = c_{1,1}$.

In this paper, we compute the metric dimension of truncated wheels.

Mathematics Subject Classification: 05C12

Keywords: resolving set, metric dimension, truncated wheel

1 Introduction

The metric dimension problem was first introduced by Harary and Melter [5] in 1976 and independently by Slater [8] in 1988. Several authors studied this topic and published numerous results. Interested readers are also referred to [1], [2], [3], [4], [6], and [9].

Let $G = (V(G), E(G))$ be a finite, simple, and connected graph. The *distance* between two vertices u and v of G, denoted by $d(u, v)$, is the length of the shortest u-v path in G. For an ordered subset $W = \{w_1, w_2, \dots, w_k\}$ of $V(G)$, we refer to the k-vector (ordered k-tuple) $r(v|W) = (d(v, w_1), d(v, w_2), d(v, w_3), \dots, d(v, w_k))$ as the *metric representation of vertex v with respect to W*. The set W is called a *resolving set of G* if every two distinct vertices u and v satisfy $r(u|W) \neq r(v|W)$. The *metric dimension of G*, denoted by $\beta(G)$, is the minimum cardinality of a resolving set of G.

For a given ordered set $W = \{w_1, w_2, , w_3 \dots, w_k\}$ of vertices of G, it is not difficult to see that the ith coordinate of $r(v|W)$ is 0 if and only if $v = w_i$. Moreover, although W is treated as an ordered set, when its elements are permuted, the coordinates of $r(v|W)$ will follow correspondingly. Thus, to show that W is a resolving set of G, it suffices to verify that $r(u|W) \neq r(v|W)$ for each pair of distinct vertices $u, v \in V(G) \backslash W$ for one particular ordering of the elements of W.

Let $n \geq 3$ be an integer. A *truncated wheel*, denoted by TW_n, is the graph with vertex set $V(TW_n) = \{a\} \cup B \cup C$, where $B = \{b_i : 1 \leq i \leq n\}$ and $C = \{c_{j,k} : 1 \leq j \leq n, 1 \leq k \leq 2\}$, and edge set $E(TW_n) = \{ab_i : 1 \leq i \leq n\} \cup \{b_i c_{i,k} : 1 \leq i \leq n, 1 \leq k \leq 2\} \cup \{c_{j,1} c_{j,2} : 1 \leq j \leq n\} \cup \{c_{j,2} c_{j+1,1} : 1 \leq j \leq n\}$, where $c_{n+1,1} = c_{1,1}$. This new graph specie was introduced by Lee [7].

In the succeeding computations, additions on the subscript of vertex b_i and the first subscript of vertex $c_{j,k}$ are taken modulo n, while that of the second subscript of $c_{j,k}$ is taken modulo 2.

In this paper, by proving several propositions and lemmas, we completely compute the metric dimension of all truncated wheels. We end the paper with a conjecture on the metric dimension of a variant of truncated wheels.

2 Main Results

Chartrand, et. al. [4] showed that, for a connected graph G of order n, $\beta(G) = 1$ if and only if G is a path of order n. It follows that $\beta(TW_n) \geq 2$ for every integer $n \geq 3$.

It is not difficult to check that $S = \{b_1, b_2\}$ and $T = \{c_{1,1}, c_{2,2}\}$ are resolving sets of TW_3 and TW_5, respectively. Thus, we have the following proposition.

Proposition 2.1. $\beta(TW_3) = \beta(TW_5) = 2$

Proposition 2.2. $\beta(TW_4) = 3$

Proof. It is not difficult to check that the set $S = \{b_1, b_2, b_3\}$ is a resolving set of TW_4. We show that there exists no resolving set S of TW_4 with order 2.

Consider an ordered subset S of $V(TW_4)$ with order 2. It suffices to consider five possibilities for S.

Case 1. If $S = \{a, b\}$ for some $b \in B$, then $r(b_k|S) = (1, 2)$ for each $b_k \neq b$ in B.

Case 2. If $S = \{b_i, b_j\}$, then $r(b_k|S) = (2, 2)$ for every $b_k \in B \backslash S$.

Case 3. If $S = \{a, c\}$ for some $c = c_{i,j}$, then $r(b_{i+j}|S) = r(b_{i+j+1}|S) = (1, 3)$.

Case 4. Suppose both vertices in S are from C.

Subcase 4.1. If $S = \{c_{i,1}, c_{i,2}\}$ or $S = \{c_{i,1}, c_{i+1,2}\}$, then $r(b_{i+3}|S) = r(c_{i+3,1}|S) = (2, 3)$.

Subcase 4.2. If $S = \{c_{i,1}, c_{i+2,2}\}$, then $r(a|S) = r(b_{i+3}|S) = (2, 2)$.

Subcase 4.3. If $S = \{c_{i,1}, c_{i+3,2}\}$, then $r(b_{i+1}|S) = r(b_{i+2}|S) = (3, 3)$.

Subcase 4.4. If $S = \{c_{i,j}, c_{i+1,j}\}$, then $r(b_{i+j+1}|S) = r(c_{i+j+1,j+1}|S) = (3, 3)$.

Subcase 4.5. If $S = \{c_{i,j}, c_{i+2,j}\}$, then $r(a|S) = r(c_{i+1,j}|S) = (2, 2)$.

Subcase 4.6. If $S = \{c_{i,j}, c_{i+3,j}\}$, then $r(b_{i+j}|S) = r(c_{i+j,j+1}|S) = (3, 3)$.

Case 5. Suppose $S = \{b, c\}$ for some $b = b_i$ and $c \in C$.

Subcase 5.1. If $c = c_{i,1}$ or $c_{i+3,2}$, then $r(b_{i+1}|S) = r(b_{i+2}|S) = (2, 3)$.

Subcase 5.2. If $c = c_{i,2}$ or $c_{i+1,1}$, then $r(b_{i+2}|S) = r(b_{i+3}|S) = (2, 3)$.

Subcase 5.3. If $c = c_{i+1,2}$, then $r(b_{i+1}|S) = r(c_{i+1,1}|S) = (2, 1)$.

Subcase 5.4. If $c = c_{i+2,1}$, then $r(b_{i+1}|S) = r(c_{i+1,1}|S) = (2, 2)$.

Subcase 5.5. If $c = c_{i+2,2}$, then $r(b_{i+3}|S) = r(c_{i+3,2}|S) = (2, 2)$.

Subcase 5.6. If $c = c_{i+3,1}$, then $r(b_{i+3}|S) = r(c_{i+3,2}|S) = (2, 1)$.

The existence of two vertices with the same metric representation with respect to S in each of the cases and subcases implies that there cannot exist a resolving set of TW_4 with order 2. Therefore, we conclude that $\beta(TW_4) = 3$. □

Lemma 2.3. *Let $n \geq 6$ be an integer. Then $\beta(TW_n) \leq \lfloor n/2 \rfloor$.*

Proof. Let $\alpha = \lfloor n/2 \rfloor$. We claim that $S = \{c_{1,1}, c_{3,1}, c_{5,1}, \ldots, c_{2\alpha-1,1}\}$ is a resolving set of TW_n. Note that $|S| = \alpha$.

We separate the cases when n is even and when n is odd. We enumerate the distinct metric representations of v with respect to S for all $v \in V(TW_n)\backslash S$.

Case 1. Suppose n is even. Then the metric representations of the vertices of TW_n are as follows.

(1.1) $r(a|S) = (2, 2, \ldots, 2)$

(1.2) For each $i = 1, 3, 5, \ldots, n-1$, $r(b_i|S)$ has $(\alpha - 1)$ 3's and one 1.

(1.3) For each $i = 2, 4, 6, \ldots, n$, $r(b_i|S)$ has $(\alpha - 1)$ 3's and one 2.

(1.4) For each $i = 2, 4, 6, \ldots, n$, $r(c_{i,1}|S)$ has $(\alpha - 2)$ 4's and two 2's.

(1.5) For each $i = 1, 3, 5, \ldots, n-1$, $r(c_{i,2}|S)$ has $(\alpha - 2)$ 4's, one 1, and one 3, where, in particular, 1 is located in the $((i+1)/2)$th coordinate.

(1.6) For each $i = 2, 4, 6, \ldots, n$, $r(c_{i,2}|S)$ has $(\alpha - 2)$ 4's, one 1, and one 3, where, in particular, 1 is located in the $((i/2) + 1)$th coordinate.

Case 2. Suppose n is odd. Then we have the following metric representations of the vertices of TW_n.

(2.1) $r(a|S) = (2, 2, \ldots, 2)$

(2.2) For each $i = 1, 3, 5, \ldots, n-2$, $r(b_i|S)$ has $(\alpha - 1)$ 3's and one 1. Moreover, we have $r(b_n|S) = (2, 3, 3, \ldots, 3)$.

(2.3) For each $i = 2, 4, 6, \ldots, n-3$, $r(b_i|S)$ has $(\alpha - 1)$ 3's and one 2. Moreover, we have $r(b_{n-1}|S) = (3, 3, \ldots, 3)$.

(2.4) $r(c_{n,1}|S) = (2, 4, 4, \ldots, 4)$

(2.5) For each $i = 2, 4, 6, \ldots, n-3$, $r(c_{i,1}|S)$ has $(\alpha - 2)$ 4's and two 2's. Moreover, we have $r(c_{n-1,1}|S) = (4, 4, \ldots, 4, 2)$.

(2.6) For each $i = 1, 3, 5, \ldots, n-4$, $r(c_{i,2}|S)$ has $(\alpha - 2)$ 4's, one 1, and one 3, where, in particular, 1 and 3 are located in the $((i+1)/2)$th and $((i+3)/2)$th coordinates, respectively. Moreover, we have $r(c_{n-2,2}|S) = (4, 4, \ldots, 4, 1)$ and $r(c_{n,2}|S) = (1, 4, 4, \ldots, 4)$.

(2.7) For each $i = 2, 4, 6, \ldots, n-3$, $r(c_{i,2}|S)$ has $(\alpha - 2)$ 4's, one 1, and one 3, where, in particular, 1 is located in the $((i/2) + 1)$th coordinate. Moreover, we have $r(c_{n-1,2}|S) = (3, 4, 4, \ldots, 4, 3)$.

It is not difficult to verify that all these metric representations with respect to S are distinct. □

The following lemma is not difficult to prove, and so we omit its proof.

Lemma 2.4. *Let $n \geq 3$. If S satisfies one of the following conditions:*

(i) $S \subset B$ with $1 \leq |S| \leq n-2$, or

(ii) $S = \{a\} \cup S_B$, where $S_B \subset B$ with $1 \leq |S_B| \leq n-2$,

then all vertices in $B \backslash S$ have the same metric representation with respect to S.

Lemma 2.5. *Let $n \geq 6$. If S satisfies one of the following conditions:*

(i) $S \subset C$ with $1 \leq |S| \leq \lfloor n/2 \rfloor - 1$, or

(ii) $S = \{a\} \cup S_C$, where $S_C \subset C$ with $1 \leq |S_C| \leq \lfloor n/2 \rfloor - 2$,

then there exist two vertices in B that are not adjacent to any vertex in $S \cap C$ and that have the same metric representation with respect to S.

Proof. We prove the result when S satisfies condition (i). The result then follows immediately when S satisfies condition (ii).

Let $\alpha = |S|$, and B' be the set of vertices in B that are not adjacent to any vertex in S. Then $|B'| \geq n - \alpha \geq 4$. We show that two vertices in B' have metric representations with respect to S that consist of α 3's.

Suppose that only at most one vertex in B' has metric representation with respect to S consisting of α 3's. Then at least $|B'| - 1$ vertices in B' have distance 2 from a vertex in S.

Since every vertex in B has exactly two vertices in C of distance 2, it follows that

$$|S| \geq |B'| - 1 \geq n - \alpha - 1 \geq n - \left(\frac{n}{2} - 1\right) - 1 = \frac{n}{2},$$

which is a contradiction. □

Lemma 2.6. *Let $n \geq 6$. If S satisfies one of the following conditions:*

(i) $S = S_B \cup S_C$, *where* $S_B \subset B$, $S_C \subset C$, $|S_B| \geq 1$, $|S_C| \geq 1$, *and* $|S| \leq \lfloor n/2 \rfloor - 1$, *or*

(ii) $S = \{a\} \cup S_B \cup S_C$, *where* S_B *and* S_C *are similarly defined as in (i), except that* $|S_B| + |S_C| \leq \lfloor n/2 \rfloor - 2$,

then there exist two vertices in $B \backslash S_B$ *with the same metric representation with respect to* S.

Proof. We prove the result when condition (i) holds. Then, because $d(b, a) = 1$ for all $b \in B$, the result immediately follows when condition (ii) holds.

Let $S' = S'_B \cup S_C$, where S'_B is the set obtained from S_B by replacing each $b_i \in S_B$ by $c_{i,1}$. Clearly, $S'_B \subset C$ and $|S'| \leq |S| \leq \lfloor n/2 \rfloor - 1$. By Lemma 2.5, there exist two vertices u and v in B that are not adjacent to any vertex in S' and that have a common metric representation with respect to S'. Because $d(b_i, b_j) = 2$ for every pair of distinct vertices b_i and b_j in B, these vertices u and v also have the same metric representation with respect to S. □

The following lemma summarizes Lemmas 2.4, 2.5, and 2.6.

Lemma 2.7. *Let* $n \geq 6$ *be an integer. Then there exists no resolving set of* TW_n *with order less than* $\lfloor n/2 \rfloor$*; that is,* $\beta(TW_n) \geq \lfloor n/2 \rfloor$.

By Propositions 2.1 and 2.2 and Lemmas 2.3 and 2.7, we have the main theorem of the paper.

Theorem 2.8. *Let* $n \geq 3$ *be an integer. Then*

$$\beta(TW_n) = \begin{cases} 2 & \text{if } n = 3 \\ 3 & \text{if } n = 4 \\ \lfloor n/2 \rfloor & \text{if } n \geq 5. \end{cases}$$

3 A Variant of TW_n

Let $n \geq 3$ be an integer. Define TW_n^* as the graph with vertex set $V(TW_n^*) = V(TW_n)$ and edge set $E(TW_n^*) = E(TW_n) \cup \{b_i b_{i+1} : 1 \leq i \leq n\}$, where $b_{n+1} = b_1$.

It is easy to check that $S = \{b_1, b_2, c_{3,1}\}$ is a resolving set of TW_3^*. It is also not difficult to show that $\beta(TW_3^*) \geq 3$. It follows that $\beta(TW_3^*) = 3$.

Moreover, it can be verified that $\beta(TW_4^*) = \beta(TW_5^*) = 2$ and $\beta(TW_6^*) = 3$.

With the initial values above, we end the paper with a conjecture on the metric dimension of TW_n^*.

Conjecture 3.1. *Let $n \geq 3$ be an integer. Then*

$$\beta(TW_n^*) = \begin{cases} 3 & \text{if } n = 3 \text{ or } n = 6 \\ 2 & \text{if } n = 4 \text{ or } n = 5 \\ \lceil n/2 \rceil - 1 & \text{if } n \geq 7. \end{cases}$$

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